Week 2 Notes

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1 Last Week's "To Think About"

Use Euclid's algorithm to compute gcd(21,13)

21=1(13)+8 13=1(8)+5 8=1(5)+3 5=1(3)+2 3=1(2)+12=2(1)

Note that the sequence of remainders is the Fibonacci Sequence backwards.

Compute $gcd(F_{n+1}, F_n)$ How many times does F_n go into F_{n+1} ?

$$F_{n+1} = 1(F_n) + F_{n-1}$$

$$F_n = 1(F_{n-1}) + F_{n-2}$$

Note: $0 \le F_{n-1} < F_n$ always holds for n > 1

The steps of Euclid's algorithm for Fibonacci numbers are: $F_{n+1} = 1(F_n) + F_{n-1}$ where q=1 if n¿1 and r=F_{n-1}

Now we can take this and write an induction proof from here.

2 Divisibility (cont.)

2.1 Corollary 2-1

If d = gcd(a, b) then there exists integers x and y such that ax + by = d.

Proof: Use the Euclidean algorithm backwards.

Suppose that we use the Euclidean algorithm and get equations
$$\begin{split} \mathbf{a}=&\mathbf{q}_0b+r_1\\ \mathbf{b}=&\mathbf{q}_1r_1+r_2\\ \mathbf{r}_{n-1}=&q_nr_n+0 \end{split}$$

Claim: There exist x_i and y_i such that $ax_i + by_i = r_i$ for each $i \ge 1$; and $ax_{i-2} + by_{i-2} = r_{i-2}$.

Proof By Induction:

Base Case: i = 1 $\mathbf{r}_1 = a - q_0 b$ Let $x_1 = 1, y_1 = -q_0$

Induction Step: Suppose there exist x_{i-1} and y_{i-1} such that $ax_{i-1} + by_{i-1} = r_{i-1}$

We want to prove there exist x_i and y_i such that $ax_i + by_i = r_i$.

We know from before that $r_{i-2} = q_{i-1}r_{i-1} + r_i$ so it follows that $r_i = r_{i-2} - q_{i-1}r_{i-1}$ = $(ax_{i-2} + by_{i-2}) - q_{i-1}(ax_{i-1} + by_{i-1})$ = $a(x_{i-2} - q_{i-1}x_{i-1}) + b(y_{i-2} - qy_{i-1})$

We can define these quantities being multiplied to a and b as x_i and y_i .

When i = n we have $ax_n + by_n = r_n = d = gcd(a, b)$.

2.2 Example

Find x and y such that 16x + 6y = gcd(16, 6) = 2

16 = 2(6) + 4 can be rewritten to 4 = (16 - 2(6))

6 = 1(4) + 2 can be rewritten to 2 = (6 - 1(4))

It follows that: 2 = 6 - 1(16 - 2(6)) 2 = 6 - 1(16) + 2(6)2 = 3(6) - 1(16)

Thus, x = -1 and y = 3.

2.3 Corollary 2-2

If gcd(a, b) = d there exist x and y such that ax + by = c if and only if d|c.

Proof: \Rightarrow Suppose ax + by = c has a solution. We know that d|a and d|b so a = df and b = dg. Then

dfx + dgy = cd(fx + gy) = c

Thus, d|c.

 \Leftarrow Suppose d|c. This means that c = de. By our theorem, there exist x and y such that ax + by = d. Multiply through ea(xe) + b(ye) = de = c

2.4 Definitions

- 1. p is a prime number if whenever a|p and a > 0 then either a = 1 and a = p. (Definition 2-2)
- 2. We say that a and b are relatively prime if gcd(a, b) = 1. (Definition 2-3)

For ex. 10 and 21 are relatively prime.

Theorem: If gcd(a, c) = 1 and a|bc then a|b

Proof: Since gcd(a, c) = 1 there exist x and y so that ax + cy = 1Multiply through by b: abx + bcy = bSince a|bc there exists f such that af = bc. So: abx + afy = ba(bx + fy) = bSo a|b

2.5 Corollary 2-3

If p is a prime number and p|ab then either p|a or p|b. (not both)

Proof: If p|a we're done. So suppose it doesn't, then gcd(a, p) = 1

By our theorem, p|b.

2.6 The Fundamental Theorem of Arithmetic

If $n \ge 2$ is an integer then there exists prime numbers $p_1 \le p_2 \le p_3 \dots \le p_k$

And this is unique n that if $p_1' \leq p_2' \leq ...p_l'$ are primes and $p_1'p_2'...p_l' = n$ then l=k and $p_i=p_i$

Proof by induction:

Base Case: n = 2; 2 is prime so we have a factorization into primes $p_1 = 2$ and is unique.

Induction: Assume that for every $1 < a < n \ a$ has a a unique factorization into primes.

Case 1: n is prime then $p_1 = n$ and we have a unique factorization of n.

Case 2: n is not prime. This means there exist a and b with 1 < a < n and 1 < b < n with n = ab.

By our induction hypothesis a and b factor uniquely into primes.

 $a = p_1 p_2 ... p_k$ $b = p'_1 p'_2 ... p'_l$ $So n = p_1 p_2 ... p_k p'_1 p'_2 ... p'_l$

Reorder to get a factorization of n.

Regarding uniqueness:

So n has a factorization, need to show that it is unique.

Suppose this factorization is not unique. Then

$$\begin{split} n &= p_1 p_2 \dots p_k = q_1 q_2 \dots q_l \\ \text{Now } p_1 | n \text{ so } p_1 | q_1 q_2 \dots q_l \\ \text{So either } p_1 | q_1 \text{ or } p_1 | (q_2 \dots q_l) \\ \text{By induction, } p_1 | q_i \text{ for some } i \text{ thus } \\ p_1 &= q_i \end{split}$$

Similarly, $q_i | p_j$ for some j and so $q_1 = p_j$ $p_1 = q_i = p_j \ge p_1$ so j = i = 1 and $p_1 = q_1$ $\frac{n}{p_1} = p_2 p_3 \dots p_k = q_2 \dots q_l$ $\frac{n}{p_1} < n$ so it factors uniquely (into prime numbers) so

k = l and $p_i = q_i$