

3/01/2018 Class Notes

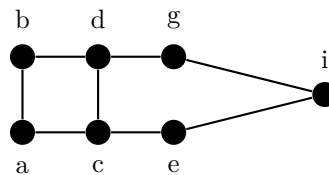
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Prim's Algorithm:

- Pick any vertex V , then pick any edge with the lowest weight connected to vertex V .
- Repeat. Taking the edge with the lowest weight that connects the existing tree to a vertex not already connected to it.

How many spanning trees does a connected graph have?

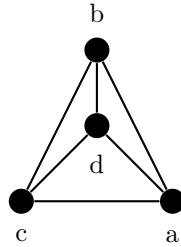
- Example:



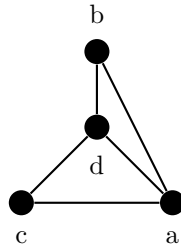
- Let the edge from d to c be represented by e_1
- First count the trees in the graph including e_1
 - * There are 3 possible edges to remove on the left side of e_1
 - * There are 4 possible edges to remove on the right side of e_1
 - * This means that including e_1 , there are 12 total possible spanning trees ($4 * 3$)
- Now count the trees in the graph not including e_1
 - * There are 7 possible edges to remove and removing any of them results in a spanning tree.
- Therefore are 19 possible spanning trees total because $12 + 7 = 19$

How many spanning trees are there on n labelled vertices if any edge is allowed?

- Counting spanning subtrees of a complete graph of order n
- Lets try this with $n = 4$

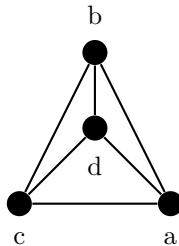


- Get rid of an edge



- Let the edge from vertex d to a be e_2
We want to count the number of spanning trees in our graph G including e_2 , as well as not including e_2 .
- Number of trees including e_2 :
2 on left * 2 on right = 4 total
Number of trees not including e_2 :
4 total
All together there are 8 in total.

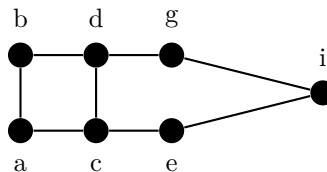
- Now lets try it including the original edge we removed (the edge from vertex b to vertex c). Call it e1.



- Now we repeat the same process as above. In our new graph, if we include e2, we have 4 possible spanning trees. If we don't include our edge e2 we have another 4 possible spanning trees. If we multiply these together we get 16 total possible spanning trees.
- The equation to calculate the number of spanning subtrees for a complete graph is n^{n-2}

Matrix Tree Theorem:

- Given a connected graph G from the matrix A where:
 $A_{ij} = deg(V_i)$
 $A_{ij} = -1$ if $V_i V_j$ exist within $E(G)$, 0 if otherwise
- Example:



- $A = \begin{bmatrix} 2 & -1 & 0 & 0 & -1 \\ -1 & 3 & -1 & -1 & 0 \\ 0 & -1 & 2 & -1 & 0 \\ 0 & -1 & -1 & 3 & -1 \\ -1 & 0 & 0 & -1 & 2 \end{bmatrix}$

- Then the number of spanning trees of our graph is equal to the determinant of the matrix $[A_{ij}]$ where this is any cofactor of the matrix obtained by removing row i and column j
- This gives us an easy way to computer the number of spanning trees even if the graph is really large.