

# Day 2 Notes

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# 1 Abstract Theorems

## 1.1 A Divides B

*Definition:* We say 'a divides b' and write  $a|b$  if  $b = ka$  for some integer  $k$ .

### 1.1.1 Examples

$$2|18 \equiv 18 = 9 \times 2$$

In this case,  $k = 9$ .

$$2|0 \equiv 0 = 0 \times 2$$

In this case,  $k = 0$ . *Note:* This means that  $a|0$  for every  $a$ .

## 1.2 B Modulo N

*Definition:*  $a \equiv b \pmod{n}$  and say "a is congruent to b modulo n" if  $n|(a - b)$ .

### 1.2.1 Examples

$$33 \equiv 3 \pmod{10}$$

$$33 - 3 = 30$$

$$10|30$$

$$27 \equiv 42 \pmod{5}$$

$$27 - 42 = -15$$

$$5|-15$$

## 1.3 Modulo Equivalency, Addition

If  $a \equiv c \pmod{n}$  and  $b \equiv d \pmod{n}$ , then  $a + b \equiv c + d \pmod{n}$ .

### 1.3.1 Examples

$12 \equiv 22 \pmod{10}$  and  $7 \equiv 37 \pmod{10}$  then...

$$12 + 7 \equiv 22 + 37 \pmod{10}$$

$$19 \equiv 59 \pmod{10}$$

### 1.3.2 Proof of Theorem

Since  $a \equiv c \pmod{n}$ , then  $n|(a-c)$  so that means  $a-c \equiv k \times n$  or  $a = c+k \times n$   
So:

$$\begin{aligned} a + b &= (c + k \times n) + (d + l \times n) \\ &= (c + d) + k \times n + l \times n \\ &= c + d + n(k + l) \end{aligned}$$

Is  $a + b \equiv c + d \pmod{n}$ ?

$$(a + b) - (c + d) \equiv -n(k + l)$$

So yes, it is congruent.

## 1.4 Module Equivalency, Multiplication

If  $a \equiv c \pmod{n}$  and  $b \equiv d \pmod{n}$ , then  $a \times b \equiv c \times d \pmod{n}$ .

### 1.4.1 Proof of Theorem

$a = c + k \times n$  and  $b = d + l \times n$ . So:

$$\begin{aligned} a \times b &= (c + k \times n) \times (d + l \times n) \\ &= cd + cln + dkn + kln^2 \\ &= cd + n \times (cl + dk + kln) \end{aligned}$$

So  $ab = cd + nj$ , where  $j$  is an integer so  $ab \equiv cd \pmod{n}$ .

## 1.5 Final Note

You can add, subtract, and multiply  $a$  and  $b$  but you can't always divide.

## 2 Affine Cipher (aka a better Caesar Cipher)

Pick a key, two integers  $\alpha$  and  $\beta$ . Assume  $0 \leq \alpha, \beta \leq 25$ . For instance,

$$E(x) = \alpha x + \beta \pmod{26}$$

### 2.1 Encryption Example

$\alpha = 7, \beta = 20$  - *This is the key.* Plaintext is 'a' or 0 19.

$$\begin{aligned} E(0) &= 7(0) + 20 \pmod{26} \\ &= 20 \end{aligned}$$

or 'U'

$$\begin{aligned} E(19) &= 7(19) + 20 \pmod{26} \\ &= 153 \pmod{26} \\ &= 23 \end{aligned}$$

or 'X'

Ciphertext is equal to 'UX'.

## 2.2 Decryption Example

$$\begin{aligned} y &= E(x) = \alpha x + \beta \pmod{26} \\ y - \beta &= \alpha x \pmod{26} \end{aligned}$$

But we can't divide all numbers (can't take the mod of a fraction). So we must find a number  $\alpha^{-1}$  where  $\alpha^{-1}\alpha \equiv 1 \pmod{26}$ . If  $\alpha^{-1}$  exists, multiply both sides by  $\alpha^{-1}$ . So:

$$\begin{aligned} \alpha^{-1}(y - \beta) &\equiv \alpha^{-1}\alpha x \pmod{26} \\ &\equiv x \pmod{26} \end{aligned}$$

Which gives us:

$$D(y) = \alpha^{-1}(y - \beta)$$

By using the table in the schedule, if  $\alpha = 7$  then  $\alpha^{-1} = 15$ .

$$\begin{aligned} D(y) &= 15(y - 20) \pmod{26} \\ &= 15y + 15 \times 6 \pmod{26} \\ &= 15y + 12 \end{aligned}$$

*Note:  $\beta$  went from -20 to 6 because we needed a positive number between 0 and 26, so we added 26 to the number.  $15 \times 6$  turned into 12 because we looked at the table to find the answer (not modulo 26).*

$$\begin{aligned} D(20) &\equiv 15(20) + 12 \pmod{26} \\ &\equiv 14 + 12 \pmod{26} \end{aligned}$$

Which is equivalent to 'a'.

$$\begin{aligned} D(23) &\equiv 15(23) + 12 \pmod{26} \\ &\equiv 7 + 12 \pmod{26} \end{aligned}$$

Which is equivalent to 't'.

### 2.2.1 Potential Problems

There can be one  $\alpha$  that can map to multiple  $\alpha^{-1}$ s. This means that the decrypted message could be wrong. What we do instead when we first pick an  $\alpha$  is make sure the row contains only one 1. This leaves out the even numbers and the number 13 from being chosen.

## 3 Attacking the Affine Cipher

### 3.1 Ciphertext Only

You can use brute force for value pairs of  $\alpha$  and  $\beta$ . This leaves 12 possibilities for  $\alpha$  and 26 for  $\beta$ , for a total of  $12 \times 26 = 312$  possibilities.

### 3.2 Chosen Plaintext

Pick  $\alpha$  to be 0, or 'a'.

$$E(0) = \alpha(0) + \beta = \beta$$

Then pick  $\alpha$  to be 1, or 'b'.

$$E(1) = \alpha(1) + \beta$$

$$E(1) - \beta = \alpha$$

### 3.3 Known Plaintext

Suppose 'it' maps to 'OH' (8, 19  $\rightarrow$  14, 7). We can use some principles in algebra to find the key  $\alpha$  and  $\beta$ .

$$\alpha(8) + \beta \equiv 14 \pmod{26}$$

$$-\alpha(19) + \beta \equiv 7 \pmod{26}$$

$$= \alpha(-11) \equiv 7 \pmod{26}$$

or

$$= \alpha(15) \equiv 7 \pmod{26}$$

We pull up the table to find  $\alpha^{-1}$  of 15, which is 7. So,

$$7 \times \alpha \times 15 \equiv 7 \times 7 \pmod{26}$$

$$\alpha \equiv 23$$

Now we plug it in to find  $\beta$ ,

$$23(8) + \beta \equiv 14 \pmod{26}$$

$$2 + \beta \equiv 14 \pmod{26}$$

$$\beta \equiv 12$$