

# MATH 314 Spring 2018 - Class Notes

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**Summary:** In Today's Class we studied into further details about the Miller Rabin Primality Test, the Factoring Trick, Dixon's Factoring Algorithm, and the examples of how each of them are implemented.

**Notes:** What is the best way to factor a very large number  $n=pq$ ?  
The first attack with just pen and paper would be trial division. We have to divide  $n$  by numbers until we find a factor.  
For example divide by 2,3,5,7,11,13 and might even have to go upto the square root. Which will take  $O(\sqrt{n})$ .

If  $n$  is an RSA modulus then,

$$n \approx 10^{240}$$

$$\sqrt{n} \approx 10^{120}$$

The second way is The Factoring Trick. If  $x^2 \equiv y^2$  and  $x \not\equiv \pm y$  then  $n$  is composite and  $d = \gcd(x - y, n)$  always going to be a non-trivial factor of  $n$ .

First attempt at Factoring Trick:

- Pick  $a$  randomly with  $\sqrt{n} < a < n$
- Lets Compare  $c = a^2 \% n$  ( $a^2$ )-is random
- If  $C = x^2$  (already a square) for some integer  $x$
- Then we have  $a^2 \equiv x^2$
- We win!!!

Pick numbers/square/reducing mod  $n$  / end on a square//

How many steps do we expect this to take?  $a$ 's are random so the  $c = a^2 \% n$  is essentially a residue mod  $n$ .

What is the probability that a random number less than  $n$  is a square?

$$\begin{aligned} & \text{number of squares less than } n / n \\ &= \sqrt{n} / n \\ \text{This probability is } & \sqrt{n} / n = 1 / \sqrt{n} \end{aligned}$$

On average we have to do this  $\sqrt{n}$  many times. So, this algorithm has running time  $O(\sqrt{n})$

### Dixon's Factoring Algorithm

Idea: Pick numbers  $a$  and computer  $c = a^2 \% n$ , keep them if  $C$  does not have any big prime factors.

Fix a bound  $B$ , prime number are considered "big" if they are bigger than  $B$ .  
We create a matrix  $F$  that has one column for every prime number, less than  $B$ .

#### • Steps:

1. Pick a random  $a\sqrt{n} < a < n$
2. Compute  $c = a^2 \% n$
3. Trial division
4. If we have not completely factored  $C$ . We give up and go back to step 1.
5. If  $C$  is completely factored into primes up to  $B$  add a row to  $F$  where the entry in each column is the number of times that prime divides  $C$ .
6. Repeat steps 1-5 until  $F$  has more rows than columns.

**Linear Algebra Fact:** Any matrix that has more rows than columns has linear dependence.

We find a linear combination of rows that we can add together to get a row where all the entries are even. Suppose we add together the rows corresponding to  $a_1, a_2, \dots, a_k$ .

$$c_1 = a_1^2 \% n, c_2 = a_2^2 \% n$$

This means that the prime factors of  $y = c_1 c_2 \dots c_k$  (All the  $C$ s multiplied together).  
All the prime numbers in  $Y$  appear to be an even power.

$$\text{so, } Y \text{ is a square } 1. \ Y = y^2 \ y = (a_1^2)(a_2^2)(a_k^2) \% n$$

$$\text{Let } x = a_1 a_2 \dots a_k$$

$$\text{Then } x^2 \equiv y^2 \pmod{n}$$