2/19/2018 Class Notes

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Definition: We say that m and n are coprime if the gcd(greatest common denominator) = 1

• Example: 10 and 21 are coprime 10 and 6 are not coprime

Theorem(Restated): The residue of a(mod m) has an inverse if and only if a and m are coprime

Chinese Remainder Theorem:

- If m and n are coprime integers then the equations $x \equiv a(modm)$ and $x \equiv b(modn)$ have a unique solution modulo m * n
- Example:
 - $x \equiv 17(mod26)$
 - $x \equiv 1 (mod2)$
 - $x \equiv 4 (mod 13)$
 - Notice that the modulo in the second two equations are coprime factors of 26.
 - The Chinese Remainder Theorem tells us that 17 is the only remainder (mod 26) that satisfies both equations
- Pick two small prime numbers
 - $x \equiv 6(mod11)$

 $x \equiv 11 (mod 13)$

- Chinese Remainder Theorem says that there is a number modulo 11*13 that satisfies both equations.
- Using these equations we can also say $x \equiv 6 + k(11)$ and using our second equation we have $x \equiv 11(mod13)$
- Therefore $6 + k(11) \equiv 11 \pmod{13}$ $k(11) \equiv 5 \pmod{13}$ $k * 11 * 6 \equiv 5 * 6 \pmod{13}$ $k \equiv 4 \pmod{13}$

- We need to find 11^{-1} using Euclids Algorithm: gcd(13,11) gives us the equation $1 \equiv 6(11) mod(13)$
- Lets try using k = 4 $x \equiv 6 + 4(11) \equiv 50$ and $50 \equiv 11 \pmod{13}$ Therefore $x \equiv 50 \pmod{143}$ is the unique solution to both $x \equiv 6 \pmod{11}$ and $x \equiv 11 \pmod{13}$

Modular Exponentiation

- Way to compute $a^k(modm)$ very fast, even if k is a very large number
- Suppose you want to compute $5^{521}(mod11)$
 - Repeated squaring lets us compute $a^{2^k} \pmod{m}$ quickly
 - Start with a and square it to get a^2 reduce (mod m), then you square it again to get $a^4,$ then again for a^8
 - Reduce (mod m) after every step to make sure the numbers don't get too big
 - Example Write k in base 2 $5^{521} = 521 = 512 + 8 + 1$ Use repeated squaring to compute $a^{2^i}(modm)$ for each power of 2 showing up in the binary for k $5^1 \equiv 5(mod11)$ $5^2 \equiv 3(mod11)$ $5^4 \equiv 9(mod11)$ $5^8 \equiv 4 \pmod{11}$ $5^{16} \equiv 5(mod11)$ $5^{32} \equiv 3(mod11)$ $5^{64} \equiv 9(mod11)$ $5^{128} \equiv 4(mod11)$ $5^{256} \equiv 5(mod11)$ $5^{512} \equiv 3(mod11)$ We know that we need the values for 2^9 , 2^3 , and 2^0 and then we multiply those values together So we have $5^{521} = 5^{521} + 5^8 + 5^0$ which by the repeated squares we did above $5^{521} = 5 * 4 * 3(mod_{11}) = 5(mod_{11})$ - This means that repeated squaring lets us compute this is time $O(\log k)$ and we never need to store numbers larger than m^2 - We can also use this to find the last digit of 3 raised to the 136th power

This is asking for $3^{136} (mod10)$ 136 = 128 + 8 $3^1 = 3(mod10)$ $3^2 = 9(mod10)$ $\begin{aligned} 3^4 &= 1(mod10) \\ 3^8 &= 1(mod10) \\ \dots \\ 3^{128} &= 1(mod10) \\ \end{aligned}$ This means that the last digit of the number would be a 1

Femat's Little Theorem

- If p is a prime number and a is coprime to p the $a^p \equiv a(modp)$ and $a^{p-1} \equiv 1(modp)$ Try this out! a=2 and p=5 $a^{p-1} \equiv 2^{5-1} \equiv 2^4 \equiv 16 \equiv 1 \pmod{5}$ Let p = 7 $a^{p-1} \equiv 2^{7-1} \equiv 2^6 \equiv 64 \equiv 1 \pmod{7}$ So it always works! - What if we try it with a non-prime number? Let p = 6 $a^{p-1} \equiv 2^{6-1} \equiv 2^5 \equiv 32 \equiv 2(mod6) \neq 1$ So it doesn't work if p is not prime! - What if we try it with base 3 $a=3\ p=5$ $a^{p-1} \equiv 3^{5-1} \equiv 3^4 \equiv 81 \equiv 1 \pmod{5}$ It works! - If we want to compute $a^k(modp)$ we can write k as some l(p-1) + rwhere r is the remainder when we divide k by p-1 $a^{k} \equiv a^{l*(p-1)} * a^{r} \equiv a^{(p-1)^{l}} * a^{r}(modp) \equiv a^{r}(modp)$ We can say this because according to Femat's little theorem $a^{(p-1)^l} \equiv$ 1 - The take home message is... We want to compute $a^{k} \pmod{p}$ if $k = r \pmod{p-1}$ then $a^{k} \equiv a^{r} \pmod{p}$ Lets try our original example again $5^{521}(mod11)$

Since $521 \equiv 1 \pmod{10}$ $5^{521} \pmod{11} \equiv 5^1 \equiv 5 \pmod{11}$

 General principle of modular arithmetic (prime version) is when you do exponentiation(mod p) you do arithmetic inside the exponent (mod p-1)