MATH 314 - Class Notes

04/18/2017

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Summary: During class, we finished SAES and began looking at the difference between Symmetric Key Cryptosystems and Public Key Crytography. We then began RSA.

Notes:

Plaintext - ARK1 - Sub - SR - MC - ARK2 - Sub - SR - ARK3 ARK1 = Key0 ARK2 = Key1 ARK3 = Key2 and Ciphertext

Decrypt SAES Receiver needs to know key Perform key expansion

- To undo ARK just ARK again
- To undo SR shift backwards
- To undo Sub Step we use inverse S-box

<u>Recall Mix Columns</u> Multiply on the right by encryption matrix To undo Mix columns we want to multiply by the inverse matrix $D = E^{-1}$ (Decryption matrix)

Ciphertext - ARK2 - SR Backwards - Inverse Sub - ARK1 - Inverse MC - SR Backwards - Inverse Sub - ARK0 ARK2 = K2 ARK1 = K1 ARK0 = K0 and plaintext

 $\begin{array}{l} \underline{AES} \\ 128 \text{ Bit Plaintext} \\ 128/192/256 \text{ bit keys} \\ (\text{Different Key Expansion}) \\ \text{Matrices we 4x4} \\ \text{work over F (one byte)} \\ (modx^8 + x^4 + x^3 + x + 1) \\ 10 \text{ rounds} \\ 9 \text{ are Sub, SR, MC, ARK} \\ \text{last is Sub, SR, ARK} \\ \text{Faster to attack AES using brute force than differential cryptanalysis (7 rounds)} \\ 3 \text{ extra round to defend against future attacks} \end{array}$

Symmetric Key Cryptosystem Alice - Encrypt - Ciphertext - Decrypt - Bob Alice and Bob have a shared secret key that is used for both encryption and decryption

Symmetric Key

- 1. AES
- 2. DES
- 3. Hill Cipher
- 4. Substitution Cipher, etc

Advantages Very Fast Very Secure if implanted right (large keys)

Disadvantages

No way to communicate with someone that doesn't already share a key with you How do you send the first message?

Public Key Cryptography

Two different keys for encryption and decryption Can't use encryption key to get decryption key Alice can publish her encryption key public Anyone can send her a message using encryption key She keeps decryption key secret She is the only one who can decrypt messages Every public key system is based on a one-way function (trapdoor function) Easy to do one way, really hard to undo

<u>RSA</u>

Trapdoor function: Multiplying integers (factoring integers) Say p and q are two 100 digit prime numbers Really easy to multiply them Compute n=pq Given n, no one knows a faster way to factor n and find p and q

Recall:

Euler's Theorem if a has gcd 1 with n then $a^{phi(n)} = 1(modn)$ When working with exponents mod n work mod phi(n) in the exponent

Suppose n = pq phi(n) = phi(p)phi(q) = (p-1)(q-1)Alice picks 2 large primes pq (120ish digits) She computes n = pqShe computes phi(n) = (p-1)(q-1)She picks an encryption exponent e need gcd(phi, phi(n)) = 1often e = 65535She computes $d = e^{-1}(modphi(n))$ Public Key (n,e) Secret Decryption Key (p,q,d) To send m to Alice, Bob computes $C = M^{(phi)}(modn)$ He sends C to Alice To decrypt Alice computes: $C = (m^{(phi(d))})(modn)$ Since ed = 1mod(phi(n)) $(m^{phi^d}) = m^{ed} = m^1(modn)$ Alice recovers Bobs message m