

Cryptography notes

James Edwards

11/28/2023

Hash Functions, not on final

-take in very large inputs to produce smaller outputs

- whenever 2 inputs produce digests, its called collision properties

1. Hard to reverse hash digest
2. Weak collision, find message that hashes to predetermine digest
3. Strong collision resistance, you create 2 messages to collide

Hash messages signature and send along side it

Elliptic Curves

Elliptic Curves are a variation to do instead of discrete log

They are the points of x and y that satisfy an equation in the form:

$$y^2 = x^3 + ax + b$$

where a and b are ints and if you draw 2 points on an elliptic curve and draw a line between them, there will be a third point that intersects the curve and the line.

1 DSA- digital Signature Algorithm

q = medium sized prime number (80 digits)

p = large prime number (200 digits)

p is chose so that $p - 1 = ql$, this means that $p - 1$ must be a multiple of q

Example.

$$q = 11, p = 67 = (6 \times 11) + 1$$

More Info

g = primitive root (mod p)

$$\alpha = g^{((p-1)/q)} = g^l \pmod{p}$$

a = private exponent, must be $0 < a < p - 1$

$$\beta = \alpha^a \pmod{p}$$

public key (p, q, α, β)

Signature Step (DSA)

Pick an ephemeral key $0 < k < p - 1$, and the $\gcd(k, p-1) = 1$

$$r = (\alpha^k \pmod{p}) \pmod{q}$$

$$s = (m + ar)k^{-1} \pmod{q}$$

(r,s) is the signature for message m so $(m, (r,s))$ is sent

Verification Step

Bob receives this message $(m,(r,s))$ and wants to check if it is valid

Bob computes:

$$U1 = s^{-1}m(modq)$$

$$U2 = s^{-1}r(modq)$$

$$\alpha^{U1}\beta^{U2}(modp)(modq)$$

This is valid if it equals r

Note, $k = s^{-1}(m + ar)(modq)$