# 9/4 Notes - Substitution and Vigenere Ciphers 

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The Caeser cipher and the Affine cipher both fall under a larger category of ciphers known as substitution ciphers.

## Substitution Cipher

## Definition

- Any encryption method that maps each individual character to either itself, or any other character, once


## Properties

- 26! combinations
- Too big to brute force
- Remains seceptible to frequency analysis

Attacks

- Chosen Plaintext
- Encrypting the entire alphabet, or a sentence containing every letter of the alphabet, will directly reveal the key.


## Examples

* "abcdefghijklmnopqrstuvwxyz"
* "The quick brown fox jumped over the lazy dog"
- Known Plaintext
* All the unique characters present in the plaintext will reveal, at least partially, the key for the cipher text
- Ciphertext Only
* frequency analysis may be used in conjunction with commonly seen patterns in language to reveal the key

Key Observation

- Simply resisting a brute force attack is not enough to gurantee the relative security of a cipher


## Vigenére Cipher

The key will consist of a word or phrase converted into their respective alphanumeric values, which we may call a vector.
Example:

- converting "key"
$-k \equiv 10 \quad(\bmod 26)$
$-e \equiv 4 \quad(\bmod 26)$
$-y \equiv 24 \quad(\bmod 26)$
- the key is $<10424>$

Encryption Steps

1. Convert the plaintext to numbers.
2. Bellow the converted plaintext, repeatedly copy the key vector until they both match in size.
3. Add the two lines in a top-down fashion then $\bmod 26$.
4. Convert the resulting line back into text.

Example

- The key will be "key" $=<10424>$
- The plaintext will be "here is how it works"
- Step 1:

| h | e | r | e | i | s | h | o | w | i | t | w | o | r | k | s |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 4 | 17 | 4 | 8 | 18 | 7 | 14 | 22 | 8 | 19 | 22 | 14 | 17 | 10 | 18 |

- Step 2 :

| h | e | r | e | i | s | h | o | w | i | t | w | o | r | k | s |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 4 | 17 | 4 | 8 | 18 | 7 | 14 | 22 | 8 | 19 | 22 | 14 | 17 | 10 | 18 |
| 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 |

- Step 3:

| h | e | r | e | i | s | h | o | w | i | t | w | o | r | k | s |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 4 | 17 | 4 | 8 | 18 | 7 | 14 | 22 | 8 | 19 | 22 | 14 | 17 | 10 | 18 |
| 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 |
| 17 | 8 | 15 | 14 | 12 | 16 | 17 | 18 | 20 | 18 | 23 | 20 | 24 | 21 | 8 | 2 |

- Step 4:

| h | e | r | e | i | s | h | o | w | i | t | w | o | r | k | s |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 4 | 17 | 4 | 8 | 18 | 7 | 14 | 22 | 8 | 19 | 22 | 14 | 17 | 10 | 18 |
| 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 | 4 | 24 | 10 |
| 17 | 8 | 15 | 14 | 12 | 16 | 17 | 18 | 20 | 18 | 23 | 20 | 24 | 21 | 8 | 2 |
| R | I | P | O | M | Q | R | S | U | S | X | U | Y | V | I | C |

- the plaintext "here is how it works" converts to "RIPO MQ RSU SX UYVIC" using the key "key"

Attacks

- Chosen Plaintext
- Having a large string of a single character, preferably 'a', will reveal the key vector
Examples
* "aаaaaaaaaaaaaaaaaaaaааааааааааааа..."
* "bbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb..."
- Known Plaintext
* Simply reverse the encryption process by subtracting the plaintext from the ciphertext
- Ciphertext Only
* Using a guessed length N for the vector key, frequency analysis may be performed for every Nth character in the ciphertext
* Babbage's Trick

1. Write the ciphertext on one line
2. copy the ciphertext on another line shifted to the right once
3. repeat step 2 until there are, at most, as many lines as the length of the ciphertext
4. Count the number of times a letter for any given line repeats when compared to the original ciphertext
Example
. We will use the previous example's cipher text "RIPOMQRSUSXUYVIC"

- Step 1:

| R | I | P | O | M | Q | R | S | U | S | X | U | Y | V | I | C |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

- Step 2/3:

| R | I | P | O | M | Q | R | S | U | S | X | U | Y | V | I | C |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| C | R | I | P | O | M | Q | R | S | U | S | X | U | Y | V | I |
| I | C | R | I | P | O | M | Q | R | S | U | S | X | U | Y | V |
| V | I | C | R | I | P | O | M | Q | R | S | U | S | X | U | Y |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\cdots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |

- step 4:

Line 1-0 coincidence(s)
Line 2-1 coincidence(s)
Line 3-2 coincidence(s)

We will see a noticeable spike in the number of coincidences on lines that are multiples of the length of the key

