

MATH 314 Fall 2019 - Class Notes

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Summary: We worked on basics principle of fermat's little theorem, 3 Pass protocol, and Euler Phi Function

Notes: when doing arithmetic (mod p). we reduce exponent (mod p-1)

Example:

$$x^3 \equiv 6 \pmod{11}$$

raise both sides to power a

$$x^{3a} \equiv 6^a \pmod{11}$$

Goal is to find a when $3a \equiv 1 \pmod{10}$

$$10 = 3 * 3 + 1$$

$$3 = 3 * 1 + 0$$

$$1 = 10 - 3 * 3$$

$$1 \equiv -3 * 3 \pmod{10}$$

$$3^{-1} \equiv 7 \pmod{10}$$

$$x^{3*7} \equiv 6^7 \pmod{11}$$

$$x^{21} = x^{10*2} \equiv 6^7 \pmod{11}$$

$$1 * x \equiv 6^7 \pmod{11}$$

$$x \equiv 6^4 * 6^2 * 6 \pmod{11}$$

Note

$$6^2 \equiv 36 \equiv 3 \pmod{11}$$

$$6^4 \equiv 3^2 \equiv 9 \pmod{11}$$

$$x \equiv 4 * 3 * 6 \pmod{11}$$

$$x \equiv 8 \pmod{11}$$

3-Pass Protocol

Physical Version

Alice want to send the mail package to bob. Alice put lock without the key to Bob. Then, Bob again put the lock without key and sends back to Alice. Alice unlock her lock and sends back to Bob. Then finally, bob unlock his key and retrieves

the message.

Math Version

Alice is going to pick big prime number p (usually 200 digits)

Alice can tell everyone about p

Then, She picks a random secret number a where $0 < a < p-1$ and $\gcd(a, p-1) = 1$

Alice's Encryption Function)

$$E(x) \equiv x^a \pmod{p}$$

Alice's Decryption Function

$$D(y) \equiv y^{a^{-1}} \pmod{p}$$

Bob also picks a secret number b where b is $0 < b < p-1$ where $\gcd(b, p-1) = 1$
he also computes b^{-1}

Note:

b^{-1} exists only if $\gcd(b, p-1) = 1$

Alice takes the plaintext and encodes it as a number m such that $0 \leq m < p$

She computes $C1 \equiv E(m) \equiv m^a \pmod{p}$

She sends $C1$ to Bob

Bob encrypts $C1$

She computes $C2 \equiv E(c1) \equiv c1^b \pmod{p}$

He sends $C2$ back to Alice

She decrypts $C2$

She computes $C3 \equiv D(c2) \equiv c2^{a^{-1}} \pmod{p}$

She sends $C3$ back to Bob

He computes $C4 \equiv D(c3) \equiv c3^{b^{-1}} \pmod{p}$

$$C4 \equiv m^{a * b * a^{-1} * b^{-1}}$$

$$C4 \equiv m \pmod{p}$$

Drawback

Have to send 3 different messages

Man in the middle attack works well against the 3- pass Protocol

Euler Phi Function

$\varphi(n)$ = no of integer a in $1 \leq a < n$ with $\gcd(a, n) = 1$

Examples:

$$\varphi(12) = 4$$

$$\varphi(26) = 12$$

$$\varphi(11) = 10$$

if 'p' is a prime then

$$\varphi(p) = p - 1$$

$$\varphi(p^k) = p^{k-1}(p - 1)$$

if the $\gcd(a, b) = 1$ $\varphi(ab) = \varphi(a) * \varphi(b)$

Examples:

$$\varphi(24) = \varphi(8 * 3)$$

$$\varphi(2^3 * 3) = \varphi(2^3) * \varphi(3)$$

$$2^2(2 - 1) * (3 - 1) = 8$$

Therefore, there are 8 numbers between 1 and 24 that doesn't have common factor with 24

$$\varphi(700) = \varphi(7 * 100)$$

$$\varphi(2^2 * 5^2) * (7) = \varphi(2^2) * \varphi(5^2) * \varphi(7)$$

$$6 * 2 * 20 = 8$$

Euler's Theorem if $\gcd(a, n) = 1$ then $a^{\varphi(n)} \equiv 1 \pmod{n}$

Basic Principle for exponents (mod n). if working in (mod n) then we work mod $\varphi(n)$ in the exponents.