10/30/19 Class Notes

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For RSA, we need big prime numbers
So we need an "is-prime" function
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fermat test(n) for k from 1 to 20:

pick a random a from (2,n-1)

compute x = a^{n-1} \pmod{n}

if(x \neq 1 \pmod{n})

return \text{ composite}

else

continue
```

return prime

The fermat test is a good way for checking primes, but it does have noticable problems

- lots of pseudoprimes
- carmichael numbers

To fix this, we will introduce the Solovay Strassen Test

```
Solovay Strassen (n):
for k from 1 to 20
pick a random a from (2,n-1)
compute x = \left[\frac{a}{n}\right] (Jacobi Symbol)
compute y = a^{(n-1)/2} \pmod{n}
```

```
if (x \not\equiv y \pmod{n})
return composite
else
continue
```

return prime

```
Example:
     Test n = 25 using a = 7
    Compute \left[\frac{a}{n}\right] \Rightarrow \left[\frac{7}{25}\right] \Rightarrow \left[\frac{25}{7}\right] \Rightarrow \left[\frac{4}{7}\right] \Rightarrow \left[\frac{2}{7}\right] \left[\frac{2}{7}\right] \Rightarrow (1)(1) \Rightarrow 1
     x = 1
     Compute y = 7^{(25-1)/2} \equiv 7^{12} \pmod{25}
     7^2 \equiv 49 \equiv 24 \pmod{25}
    7^4 \equiv (7^2)^2 \equiv 24^2 \equiv (-1)^2 \equiv 1 \pmod{25}
     7^8 \equiv 1 \pmod{25}
     7^{12} \equiv 1 \pmod{25}
     Solovay Strassen says "probably prime" because x\equiv y
          Now pick a new a
     a = 3
     Compute \left[\frac{a}{n}\right] \Rightarrow \left[\frac{3}{25}\right] \Rightarrow \left[\frac{25}{3}\right] \Rightarrow \left[\frac{1}{3}\right] \Rightarrow 1
     Compute 3^{(25-1)/2}
     3^{12} \equiv 3^8 \times 3^4
     3^2 \equiv 9 \pmod{25}
    3^{4} \equiv 9^{2} \equiv 81 \equiv 6 \pmod{25}

3^{8} \equiv (3^{4})^{2} \equiv 6^{2} \equiv 36 \equiv 11 \pmod{25}

3^{8} \times 3^{4} \equiv 6 \times 11 \equiv 66 \equiv 16 \pmod{25}
     Solovay Strassen says 25 is composite because 1 \not\equiv 16
There is a test that's even better than Solovay Strassen, with fewer
pseudoprimes
Miller-Rabin test(n):
    Write 2^k m [m is odd]
    Pick a random a in [2, n-1] (the goal is to compute a^{n-1} \pmod{n})
     Step 1: compute b_0 \equiv a^m \pmod{n}
          If b_0 \equiv 1 \text{ or } -1 \pmod{n}
               n is probably prime (continue)
     Step 2:
          for j from 1 to k-1 [k is the power from 2]
               compute b_j = b_{j-1}^2 \pmod{n}
               if b_j \equiv 1 \pmod{n}
                   return composite
               if b_j \equiv -1 \pmod{n}
```

say probably prime [continue]

If we finish the loop and $b_{k-1} \not\equiv \pm 1 \pmod{n}$ return composite Else the number is prime

For Miller-Rabin, at most, 1/4 of the possible a values say $\ensuremath{\text{prime}}$ for composite n